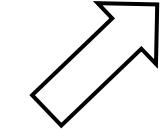


Review: Lab03 & Hw03



今日实验课内容

这里有页码
提问请指出第几页



- 作业题选讲：
 - lab03p5: Maximum Subsequence
 - hw03p2: Ping-pong
 - hw03p4: Count Change
 - hw03p6: Multiadder
 - hw03p8: All-Ys Has Been
- 一些学习方面的建议



不想听的话可以直接开搞lab04了

lab03p5: Maximum Subsequence

- `max_subseq(n, l)`: 求n的长度不超过l子序列构成的最大的数

$$\text{max_subseq}(20125, 3) = 225$$

- Base cases:



lab03p5: Maximum Subsequence

- `max_subseq(n, l)`: 求n的长度不超过l子序列构成的最大的数

$$\text{max_subseq}(20125, 3) = 225$$

- *Base cases:*
 - $n == 0$: 数字是0, 结果只可能是0
 - $l == 0$: 长度为0的子序列只可能是0
- 注意: 不要选 $n < 10$ 或者 $l == 1$ 作为base case, 会变得很复杂

lab03p5: Maximum Subsequence

- `max_subseq(n, l)`: 求n的长度不超过l子序列构成的最大的数

$$\text{max_subseq}(20125, 3) = 225$$

- $n > 0$ and $l > 0$?

- 只需要考虑末尾一位在不在答案里就够了
- 不在答案: `return max_subseq(n // 10, l)`
- 在答案里:



lab03p5: Maximum Subsequence

- `max_subseq(n, l)`: 求n的长度不超过l子序列构成的最大的数

$$\text{max_subseq}(20125, 3) = 225$$

- $n > 0$ and $l > 0$?
 - 只需要考虑末尾一位在不在答案里就够了
 - 不在答案: `return max_subseq(n // 10, l)`
 - 在答案里: `return max_subseq(n // 10, l - 1) * 10 + (n % 10)`

lab03p5: Maximum Subsequence

- `max_subseq(n, l)`: 求n的长度不超过l子序列构成的最大的数

$$\text{max_subseq}(20125, 3) = 225$$

- $n > 0$ and $l > 0$?

```
return max(max_subseq(n // 10, l),  
          max_subseq(n // 10, l - 1) * 10 + (n % 10))
```

hw03p2: Ping-pong

- 求Pingpong数列的第n项
- 用while写: SO EASY!
- 用递归写:



hw03p2: Ping-pong

- 求Pingpong数列的第n项
- 一种错误的想法：
 - Base case: $n == 1 \Rightarrow 1$
 - Recursion: $\text{pingpong}(n) = \text{pingpong}(n-1) +1/-1$
- 整了半天，做不出来，或者程序跑得贼慢 ?

hw03p2: Ping-pong

- 求Pingpong数列的第n项
- 换个思路：先写while的做法，然后改成递归（由莉莉丝提供代码）

```
1 def pingpong(n):
2     curr, result, direct = 1, 1, 1
3     while curr != n:
4         if curr % 6 == 0 or number_of_six(curr) > 0:
5             result = result - direct
6             direct = -direct
7         else:
8             result = result + direct
9             # direct = direct
10            curr = curr + 1
11    return result
```

hw03p2: Ping-pong

- 求Pingpong数列的第n项
- 换个思路：先写while的做法，然后改成递归（由莉莉丝提供代码）

```
1 def pingpong(n):
2     def helper(curr, result, direct):
3         if curr == n:
4             return result
5         if curr % 6 == 0 or number_of_six(curr) > 0:
6             return helper(curr + 1, result - direct, -direct)
7             return helper(curr + 1, result + direct, direct)
8     return helper(1, 1, 1)
```

hw03p2: Ping-pong

- 求Pingpong数列的第n项

- 再换个思路：

先确定n的时候的方向
然后递归求解
(由莉莉丝提供代码)

```
1 def pingpong(n):
2     def ping_pong(k, f):
3         if k <= 6:
4             return k
5         elif k % 6 == 0 or number_of_six(k) > 0:
6             return ping_pong(k - 1, -f) - f
7         else:
8             return ping_pong(k - 1, f) + f
9     def add_or_sub(k):
10        if k < 6:
11            return 1
12        elif k % 6 == 0 or number_of_six(k) > 0:
13            return -1 * add_or_sub(k - 1)
14        else:
15            return add_or_sub(k - 1)
16    return pingpong(n, add_or_sub(n))
```

hw03p2: Ping-pong

- 求Pingpong数列的第n项
- 再换个思路： $\text{pingpong}(k)$ 和 $\text{pingpong}(n)$ 相差多少呢？

hw03p2: Ping-pong

- 求Pingpong数列的第n项
- 再换个思路： pingpong(k)和pingpong(n)相差多少呢？

- Base case: $k == n \Rightarrow ?$

hw03p2: Ping-pong

- 求Pingpong数列的第n项
- 再换个思路： pingpong(k)和pingpong(n)相差多少呢？
 - Base case: $k == n \Rightarrow 0$
 - Recursion: $\text{diff}(k) = ?$

hw03p2: Ping-pong

- 求Pingpong数列的第n项
- 再换个思路： $\text{pingpong}(k)$ 和 $\text{pingpong}(n)$ 相差多少呢？
 - Base case: $k == n \Rightarrow 0$
 - Recursion:
$$\begin{aligned} \text{diff}(k) &= \text{pingpong}(n) - \text{pingpong}(k) \\ &= \text{pingpong}(n) - (\text{pingpong}(k+1) - (1/-1)) \\ &= \text{diff}(k + 1) + (1/-1) \end{aligned}$$

hw03p2: Ping-pong

- 求Pingpong数列的第n项
- 再换个思路： pingpong(k)和pingpong(n)相差多少呢？
 - Base case: $k == n \Rightarrow 0$
 - Recursion: $\text{diff}(k) = \text{diff}(k + 1) + (1/-1)$



加1还是-1怎么算？

hw03p2: Ping-pong

- 求Pingpong数列的第n项
- 再换个思路： pingpong(k)和pingpong(n)相差多少呢？
 - Base case: $k == n \Rightarrow 0$
 - Recursion: $\text{diff}(k, d) = \text{diff}(k + 1, -d \text{ if } \dots \text{ else } d) + d$
 - Answer: $\text{pingpong}(n) = \text{diff}(0, 1)$

hw03p4: Count Change

- `count_change(total, next_money)`: 求找零的方法数量
- Base case:
 - $\text{total} < 0 \Rightarrow$?
 - $\text{total} == 0 \Rightarrow$?
 - $\text{money is None} \Rightarrow$?

hw03p4: Count Change

- `count_change(total, next_money)`: 求找零的方法数量
- Base case:
 - $\text{total} < 0 \Rightarrow 0$
 - $\text{total} == 0 \Rightarrow 1$
 - $\text{money is None} \Rightarrow 0$
- Recursion:
 - 多用一张当前的钱: `count(total - money, money)`
 - 用下个更大的面额: `count(total, next_money(money))`

hw03p6: Multiadder

- `multiadder(n)`: 返回一个可以连续调用`n`次的求和函数



hw03p6: Multiadder

- `multiadder(n)`: 返回一个可以连续调用n次的求和函数
- 实现 `lambda x1: lambda x2: ...: lambda xn: x1 + x2 + ... + xn`

hw03p6: Multiadder

- `multiadder(n)`: 返回一个可以连续调用`n`次的求和函数
- 实现 `lambda x1: lambda x2: ...: lambda xn: x1 + x2 + ... + xn`
- Base case: 

hw03p6: Multiadder

- `multiadder(n)`: 返回一个可以连续调用`n`次的求和函数
- 实现 `lambda x1: lambda x2: ...: lambda xn: x1 + x2 + ... + xn`
- Base case: `n == 1 => lambda x: x`

hw03p6: Multiadder

- `multiadder(n)`: 返回一个可以连续调用n次的求和函数
- 实现 `lambda x1: lambda x2: ...: lambda xn: x1 + x2 + ... + xn`
- Base case: $n == 1 \Rightarrow \text{lambda } x: x$
- Recursion: 

hw03p6: Multiadder

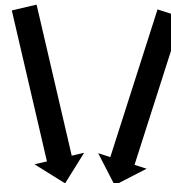
- `multiadder(n)`: 返回一个可以连续调用`n`次的求和函数
- 实现 `lambda x1: lambda x2: ...: lambda xn: x1 + x2 + ... + xn`



`lambda x12: lambda x3: ...: lambda xn: x12 + x3 + ... + xn`

hw03p6: Multiadder

- `multiadder(n)`: 返回一个可以连续调用n次的求和函数
- 实现 `lambda x1: lambda x2: ...: lambda xn: x1 + x2 + ... + xn`



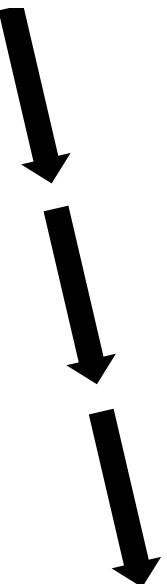
`lambda x12: lambda x3: ...: lambda xn: x12 + x3 + ... + xn`



`lambda x123: lambda x4: ...: lambda xn: x123 + x4 + ... + xn`

hw03p6: Multiadder

- `multiadder(n)`: 返回一个可以连续调用`n`次的求和函数
- `lambda x123: lambda x4: ...: lambda xn: x123 + x4 + ... + xn`

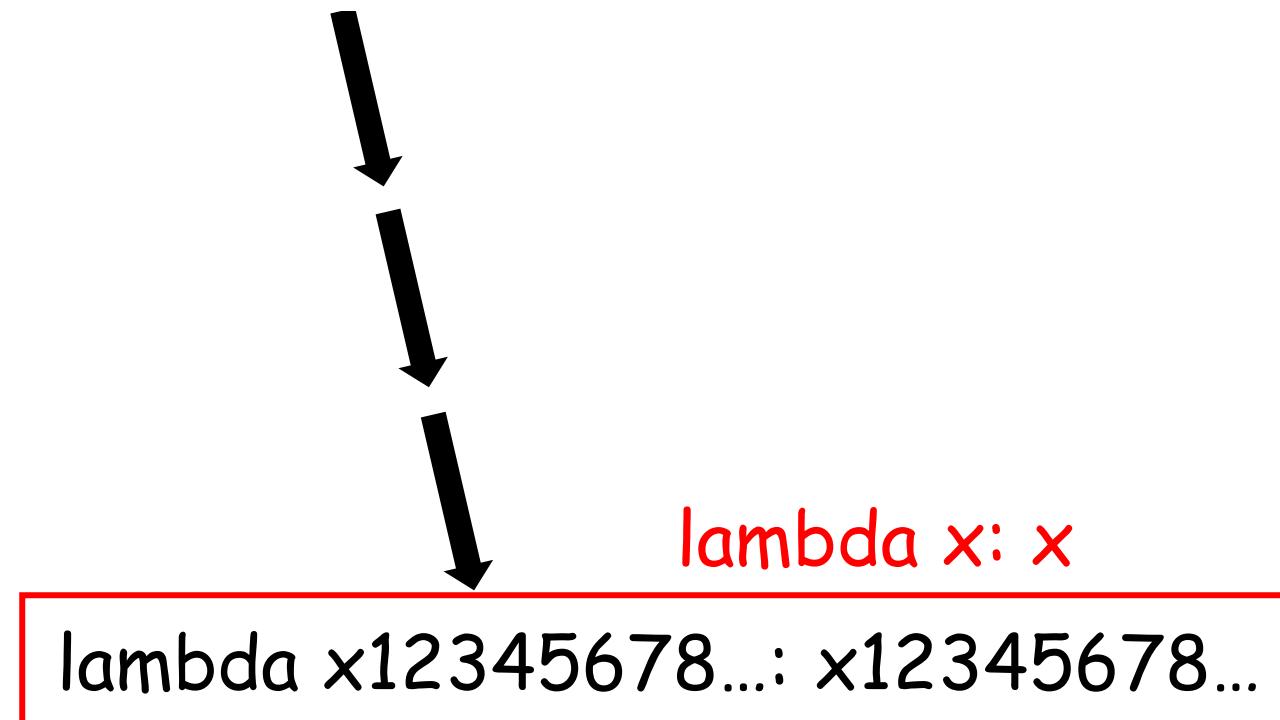


好像哪里见过?

`lambda x12345678...: x12345678...`

hw03p6: Multiadder

- `multiadder(n)`: 返回一个可以连续调用`n`次的求和函数
- `lambda x123: lambda x4: ...: lambda xn: x123 + x4 + ... + xn`



hw03p6: Multiadder

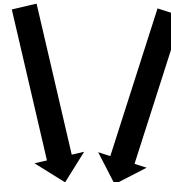
- `multiadder(n)`: 返回一个可以连续调用`n`次的求和函数
- 实现 `lambda x1: lambda x2: ...: lambda xn: x1 + x2 + ... + xn`



`lambda x12: lambda x3: ...: lambda xn: x12 + x3 + ... + xn` 

hw03p6: Multiadder

- `multiadder(n)`: 返回一个可以连续调用`n`次的求和函数
- 实现 `lambda x1: lambda x2: ...: lambda xn: x1 + x2 + ... + xn`



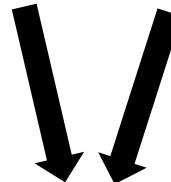
```
lambda x12: lambda x3: ...: lambda xn: x12 + x3 + ... + xn
```

multiadder(`n - 1`)



hw03p6: Multiadder

- `multiadder(n)`: 返回一个可以连续调用n次的求和函数
- 实现 `lambda x1: lambda x2: ...: lambda xn: $x_1 + x_2 + \dots + x_n$`



```
lambda x1: lambda x2: ...: lambda xn:  $x_1 + x_2 + \dots + x_n$ 
```

`multiadder(n - 1)(x1 + x2)`



!!!!!!

hw03p6: Multiadder

- `multiadder(n)`: 返回一个可以连续调用n次的求和函数
- 实现 `lambda x1: lambda x2: ...: lambda xn: x1 + x2 + ... + xn`
- Base case: $n == 1 \Rightarrow \text{lambda } x: x$
- Recursion: 

hw03p6: Multiadder

- `multiadder(n)`: 返回一个可以连续调用`n`次的求和函数
- 实现 `lambda x1: lambda x2: ...: lambda xn: x1 + x2 + ... + xn`
- Base case: `n == 1 => lambda x: x`
- Recursion: `lambda x: lambda y: multiadder(n - 1)(x + y)`



hw03p6: Multiadder

- `multiadder(n)`: 返回一个可以连续调用`n`次的求和函数
- 换个思路: E.g. `multiadder(5)(1)(2)`
 - 总共有5个数要求和, 已经运算过了2个数
 - 还要加3个数, 已有的数的和为3
 - `multiadder`的参数不够保存这些信息, 需要我们定义新的函数

hw03p6: Multiadder

- `multiadder(n)`: 返回一个可以连续调用n次的求和函数
- 换个思路: (由莉莉丝提供代码)

```
1 def multiadder(n):
2     def helper(n, curr_sum):
3         def inner(x):
4             if n == 1:
5                 return curr_sum + x
6                 return helper(n - 1, curr_sum + x)
7         return inner
8     return helper(n, 0)
```

hw03p8: All-Ys Has Been



hw03p8: All-Ys Has Been



Y是个什么玩意?

- 函数的不动点(fix-point)

$$\{ x \mid f(x) = x \}$$

hw03p8: All-Ys Has Been



Y是个什么玩意?

- 函数的不动点(fix-point)

$$\{ x \mid f(x) = x \}$$

- 不动点组合子(fix-point combinator)

$$\forall f, Y(f) = f(Y(f))$$

对于任意函数, Y可以找到他的不动点!

hw03p8: All-Ys Has Been



Y是个什么玩意?

- Y组合子(Y combinator)

$$Y = \lambda f. (\lambda x. f(x\ x))(\lambda x. f(x\ x))$$

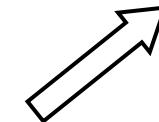
Haskell Brooks Curry was an American mathematician and logician. Curry is best known for his work in combinatory logic. While the initial concept of combinatory logic was based on a single paper by Moses Schönfinkel, Curry did much of the development. Curry is a



hw03p8: All-Ys Has Been

- Θ组合子(Θ combinator)

$$\Theta = (\lambda x. \lambda y. y(x x y))(\lambda x. \lambda y. y(x x y))$$



这是hw03p7的一种答案

Alan Mathison Turing was an English mathematician, computer scientist, logician, cryptanalyst, philosopher, and theoretical biologist. Turing was highly influential in the development of theoretical computer science, providing a formalisation of the concept



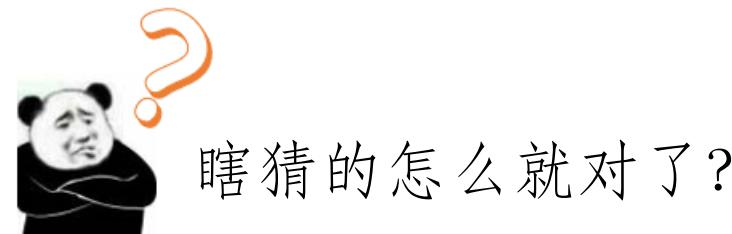
hw03p8: All-Ys Has Been

- 已知 $\lambda \text{ fib} = \text{fib } (\lambda \text{ fib})$
- fib 应该填什么能得到斐波那契数列?



hw03p8: All-Ys Has Been

- 已知 $\lambda y \text{ fib} = \text{fib } (\lambda y \text{ fib})$
- fib 应该填什么能得到斐波那契数列?
- $\text{fib} = \lambda f: \lambda r: 1 \text{ if } r \leq 1 \text{ else } f(r - 1) + f(r - 2)$



hw03p8: All-Ys Has Been

- 已知 $\text{Y fib} = \text{fib}(\text{Y fib})$
- fib 应该填什么能得到斐波那契数列?
- $\text{fib} = \lambda f: \lambda r: 1 \text{ if } r \leq 1 \text{ else } f(r - 1) + f(r - 2)$

$\text{Y fib} = \text{fib}(\text{Y fib})$

$= \lambda r: 1 \text{ if } r \leq 1 \text{ else } (\text{Y fib})(r - 1) + (\text{Y fib})(r - 2)$

- $\text{Y fib} = ?$

hw03p8: All-Ys Has Been

What Would Python Display?

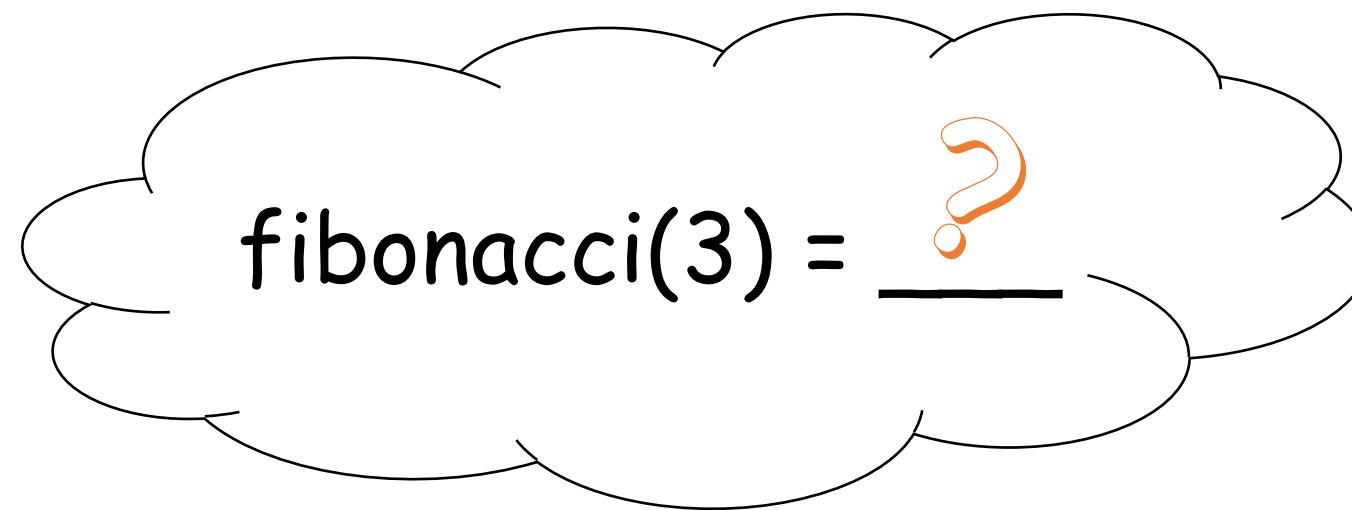
```
>>> Y = lambda f: (lambda x: f(x(x)))(lambda x: f(x(x)))
```

```
>>> Y(fib)(3) ?
```

hw03p8: All-Ys Has Been

What Would Python Display?

```
>>> Y = lambda f: (lambda x: f(x(x)))(lambda x: f(x(x)))  
>>> Y(fib)(3)
```



hw03p8: All-Ys Has Been

What Would Python Display?

```
>>> Y = lambda f: (lambda x: f(x(x)))(lambda x: f(x(x)))  
>>> Y(fib)(3)
```

Runtime Error: maximum recursion depth exceeded



hw03p8: All-Ys Has Been

What Would Python Display?

```
>>> Y = lambda f: (lambda x: f(x(x)))(lambda x: f(x(x)))  
>>> Y(fib)(3)
```

Runtime Error: maximum recursion depth exceeded

$$\begin{aligned}Y(f) &= (\lambda x: f(x(x)))(\lambda x: f(x(x))) \\&= f((\lambda x: f(x(x)))(\lambda x: f(x(x)))) \\&= f(Y f)\end{aligned}$$

hw03p8: All-Ys Has Been

What Would Python Display?

```
>>> Y = lambda f: (lambda x: f(x(x)))(lambda x: f(x(x)))  
>>> Y(fib)(3)
```

Runtime Error: maximum recursion depth exceeded

$$\begin{aligned}Y(f) &= (\lambda x: f(x(x)))(\lambda x: f(x(x))) \\&= f((\lambda x: f(x(x)))(\lambda x: f(x(x)))) \\&= f(Y\ f) = f(f(Y\ f)) = f(f(f(Y\ f))) = f(f(f(f(Y\ f)))) \\&= \dots \Rightarrow \text{Runtime Error}\end{aligned}$$

λ s hw03p8: All-~~X~~ Has Been

- Z 组合子(Z combinator)

$$Z = \lambda f. (\lambda x. f(\lambda z. x x z))(\lambda x. f(\lambda z. x x z))$$
$$\forall f, Z(f)(z) = f(Z(f))(z)$$

λ s hw03p8: All-~~X~~ Has Been

- Z 组合子(Z combinator)

$$Z = \lambda f. (\lambda x. f(\lambda z. x x z))(\lambda x. f(\lambda z. x x z))$$

$$\forall f, Z(f)(z) = f(Z(f))(z)$$

通过添加变量 z , 巧妙地避免了 λ 组合子求值时的无穷递归

$$\begin{aligned} Z(f) &= (\lambda x. f(\lambda z. x(x)(z)))(\lambda x. f(\lambda z. x(x)(z))) \\ &= f(\lambda z. (\lambda x. f(\lambda w. x(x)(w)))(\lambda x. f(\lambda w. x(x)(w))))(z) \\ &= f(\lambda z. Z(f)(z)) \end{aligned}$$

Zs

hw03p8: All-~~X~~ Has Been

- Z组合子(Z combinator)

$$Z = \lambda f. (\lambda x. f(\lambda z. x x z))(\lambda x. f(\lambda z. x x z))$$

$$\forall f, Z(f)(z) = f(Z(f))(z)$$

通过添加变量z，巧妙地避免了Y组合子求值时的无穷递归

$$\begin{aligned} Z(f) &= (\lambda x. f(\lambda z. x(x)(z)))(\lambda x. f(\lambda z. x(x)(z))) \\ &= f(\lambda z. (\lambda x. f(\lambda w. x(x)(w)))(\lambda x. f(\lambda w. x(x)(w))))(z) \\ &= f(\lambda z. Z(f)(z)) \end{aligned}$$

这和Y(f)有什么区别 ?

Z_s

hw03p8: All-~~X~~ Has Been

- Z 组合子(Z combinator)

$$Z = \lambda f. (\lambda x. f(\lambda z. x x z))(\lambda x. f(\lambda z. x x z))$$

$$\forall f, Z(f)(z) = f(Z(f))(z)$$

通过添加变量 z , 巧妙地避免了 λ 组合子求值时的无穷递归

$$Z(f) = f(\lambda z. Z(f)(z))$$

$$Z \text{ fib} = \lambda r. 1 \text{ if } r \leq 1 \text{ else } (\lambda z. Z(\text{fib})(z))(r - 1) + (\lambda z. Z(\text{fib})(z))(r - 2)$$

Z_s

hw03p8: All-~~X~~ Has Been

- Z 组合子(Z combinator)

$$Z = \lambda f. (\lambda x. f(\lambda z. x x z))(\lambda x. f(\lambda z. x x z))$$

$$\forall f, Z(f)(z) = f(Z(f))(z)$$

通过添加变量 z , 巧妙地避免了 λ 组合子求值时的无穷递归

$$Z(f) = f(\lambda z. Z(f)(z))$$

$$Z \text{ fib} = \lambda r. 1 \text{ if } r \leq 1 \text{ else } (Z \text{ fib})(r - 1) + (Z \text{ fib})(r - 2)$$

一些学习建议

不要找助教私聊问问题（聊天、心态崩了除外）

- 群里人多，并且大概率有人已经问过了
- 你的问题可能过于弱智，助教也不知道如何回答是好

- 为什么不是满分 → **亲亲，您的代码有bug哦~**
- 报错是怎么回事 → 看不懂英语的话可以买一本词典

遇到问题：

```
Traceback (most recent call last):
  return █
          ^
SyntaxError: invalid syntax
```

遇到问题：

```
Traceback (most recent call last):
  return █
          ^
SyntaxError: invalid syntax
```

invalid 在英语-中文 (繁体) 词典中的翻译



←查字典：0元，还能学英语

invalid

adjective

UK 听 /ɪn'vælɪd/ US 听 /ɪn'vælɪd/



An **invalid** document, ticket, law, etc. is not legally or officially acceptable.

(文件、票、法律等)無效的，不合法的，官方不承認的

syntax 在英语-中文 (繁体) 词典中的翻译



syntax

noun [U] • LANGUAGE • specialized

UK 听 /'sɪn.tæks/ US 听 /'sɪn.tæks/



the grammatical arrangement of words in a sentence

句法；句子結構

遇到问题：

```
Traceback (most recent call last):
  return [REDACTED]
^
SyntaxError: invalid syntax
```

invalid 在英语-中文 (繁体) 词典中的翻译

invalid

adjective

UK /ɪn'vælɪd/ US /ɪn'vælɪd/



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the grammatical arrangement of words in a sentence

句法；句子結構

←查字典：0元，还能学英语

↓群内提问：建议收费500元
知识付费 明码标价不坑人



一些学习建议

不要死磕某一道题

- 做不出就是做不出，可能是思路不对，可能是你没学会
- 一个小时都做不出的话，建议呼吸一下新鲜空气
- 再仔细读一遍题目、再看一遍ppt、录屏
- 找会做的同学/舍友问问他们的思路（如何提问.pdf）
- 放弃这道题，不要因为某一题影响整门课和其他课的学习

Questions?

201220195 叶恒迪 2021/10/25 1:12:31

恭喜你变得更强了 

201220098 杨林 2021/10/25 1:12:19

恭喜你变得更强了 

201870214 沈珺妍 2021/10/25 1:05:56

恭喜你变得更强了 

211220184 蒋知睿 2021/10/25 1:03:52

恭喜你变得更强了 

211108100-王一安 2021/10/25 1:03:38

恭喜你变得更强了 

215220005 翁邱一洪 2021/10/25 0:56:40

恭喜你变得更强了 

211220177 郑恒彬 😊 2021/10/25 0:55:04

恭喜你变得更强了 

201220195 叶恒迪 2021/10/25 0:53:55

恭喜你变的更强了 

201870214 沈珺妍 2021/10/25 0:53:34

恭喜你变得更强了 